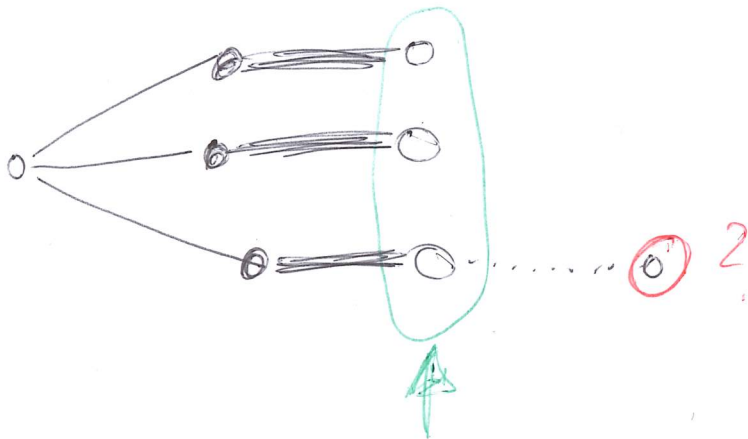
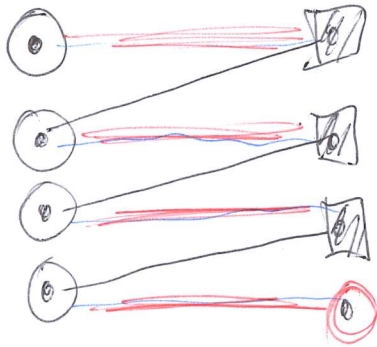


ad Matching

slide 206:



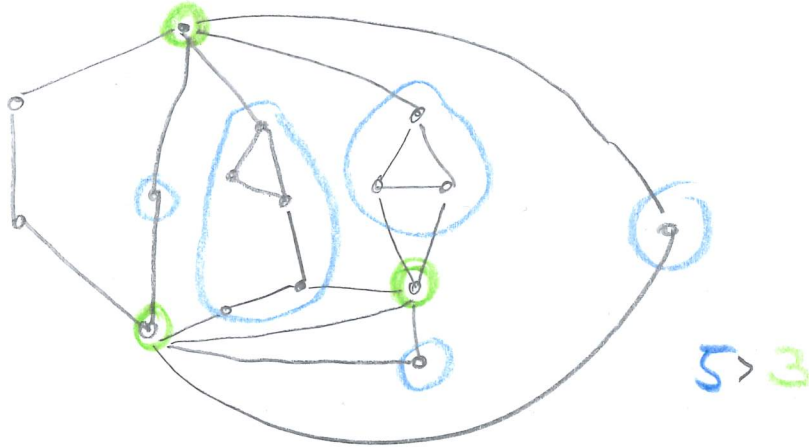
Consider: searching for perfect matching in general graphs?

Given: $G = (V, E)$

Output: Existence a perfect matching in G ?

YES \rightarrow give matching

NO \rightarrow Th. 9.16 set of vertices splitting graph in more odd components.



Theorem 9.16 (Tutte 1947)

A graph $G = (V, E)$ has a perfect matching

\Leftrightarrow For each set $A \subseteq V$: $o(G \setminus A) \leq |A|$

More general:

Theorem 9.17 (Tutte - Berge 1958) ^{formula}

For $G = (V, E)$

$\max \{ |M| \mid M \text{ is matching} \} = \min \left\{ \frac{1}{2} (|V| - o(G \setminus A) + |A|) \mid A \subseteq V \right\}$

\hookrightarrow set A gives ^{upper} bound for card. of a matching

In the algorithm for bipartite matching?

Theorem 9.18: Let G be bipartite, M a matching, T an alternating tree, for which no edge of G connects a vertex in $V(T)$ with a vertex not in $V(T)$. Then G does not have a perf. matching

proof: Consider $B(T) := V(T) \setminus W(T)$ - the set of black vertices in T



By construction: $|B(T)| = |W(T)| - 1$



G bipartite \rightarrow no two white vertices adjacent

\Rightarrow Each vertex in $B(T)$ is an odd component of $G \setminus B(T)$

$|B(T)| < |W(T)|$

\Rightarrow statement holds. \square

For general graphs?

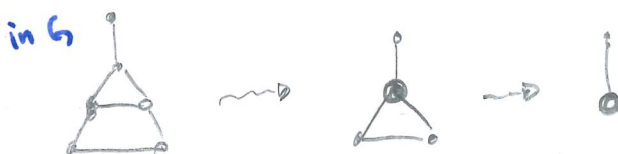
start with graph G

new graph G' (obtained from G by a series of shrinkings of odd cycles)

vertices in G' :

- "original vertices"

- "pseudo-vertices": represent shrunk odd cycles



\Rightarrow Each vertex in $V(G')$ represents an odd set of vertices

$S(v) \subseteq V(G)$

\uparrow union of odd sets

Theorem 9.19: Let G' be derived graph from G , M' matching in G' , T' ~~alternating tree~~ M' -alternating tree of G' , such that no element of $B(T)$ (black) is a pseudo vertex.

This is frustrated!

If each edge has with one end in $W(T)$ has other end in $B(T)$, then G has no perfect matching.

proof: Delete $B(T)$ from G . In G each $v \in W(T)$ is an odd component in S_v

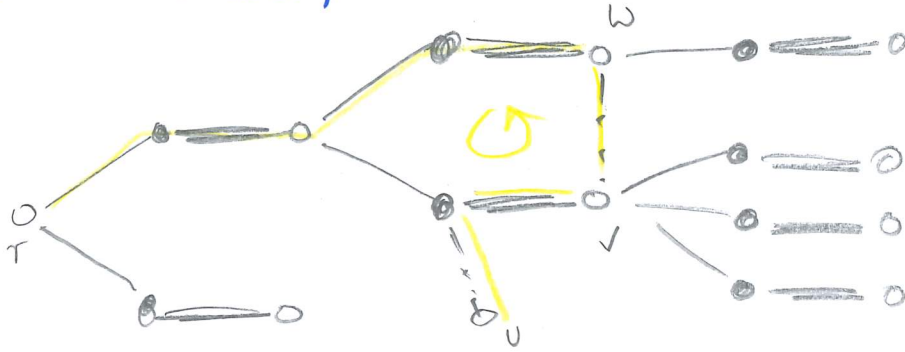
$\Rightarrow oc(G \setminus B(T)) > |B(T)| \Rightarrow$ no perfect matching \square

~~\Rightarrow ~~stages~~~~

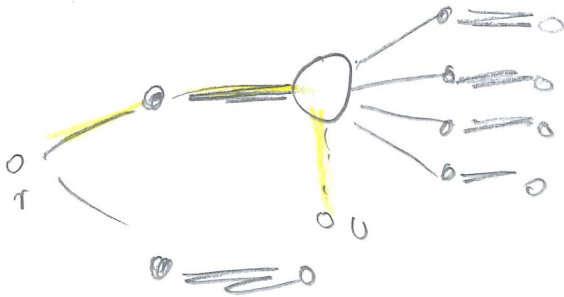
- ⇒ • derived graph has perf. matching ⇒ obtain perf. matching in G (III)
- —//— gives contradiction ⇒ no —//—//—

How to find cycles to shrink?

- edges WCT) to WCT)



after shrinking



⇒ slides

End of Alg. 6.20:

- perf. matching

- a frustrated tree T :
- ~~the~~ root unmatched
- deleting black vertices shows that no matching partner exists for all white vertices

Done?



No, can have unmatched vertices elsewhere

↳ continue search there

⇒ slides (211)

proof 9.23:

